

SHAMBHUNATH INSTITUTE OF ENGINEERING AND TECHNOLOGY

Subject :Operating System

Subject Code :KCS-401

B.Tech. : 2nd Year SEMESTER : 4th

SOLUTION FIRST SESSIONAL EXAMINATION, EVEN SEMESTER, (2019-2020)

Branch : Computer Science and Engineering

Time – 1 hr. 30 min.

Max Marks – 30

SECTION – A

Question 1(a): short-term scheduler (CPU scheduler)

Question 1(b): Output is 15

Question 1(c): SJF scheduling algorithm

Question 1(d): Round-robin scheduling

Question 1(e): 11

SECTION – B

Question 2(a):

Long Term Scheduler/Job Scheduler

- ❖ It selects process from Job Queue and assigns it to Ready Queue.
- ❖ It changes state of processes from New State to Ready State.

Short Term Scheduler/CPU Scheduler

- ❖ It selects process from Ready Queue and assigns it to CPU.
- ❖ It changes state of processes from Ready State to Running State.
- ❖ Dispatcher is responsible for saving the context of one process (i.e. content of PCB) and loading the context of another process.

Medium Term Scheduler

- ❖ It removes the processes from Main Memory.
- ❖ It reduces degree of multiprogramming.
- ❖ It is responsible for swapped out process.

Dispatcher: Dispatcher is part of Short Term Scheduler that performs following functions.

- ❖ Context switching
- ❖ Switching to User Mode
- ❖ Jumping to the proper location in the user program to restart that program

Question 2(b):

Life Cycle of Process

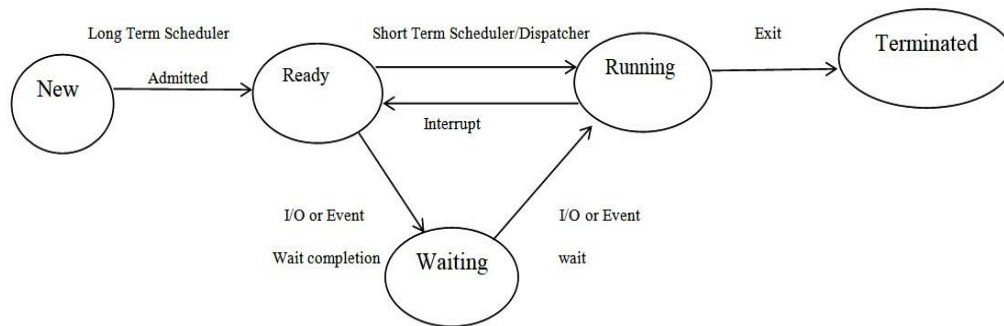
New State: When a process is created it is in new state. In new state process resides in Job Queue (Secondary Memory).

Ready State: State of process is called Ready State when Process resides in Ready Queue (Primary Memory) and waiting for CPU.

Running State: State of process is called Running state if CPU is assigned to process.

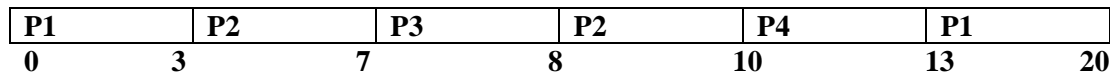
Waiting State: State of process is called Waiting State when Process resides in waiting Queue (Primary Memory) and waiting for some event like I/O to occur.

Terminated State: State of process is called Terminated State when the Process has completed its execution successfully.



Question 2(c):

Gantt Chart:



Process TAT=Completion time – Arrival time

- P1 20
- P2 7
- P3 1
- P4 5

Avg TAT= 33/4=8.25

WT:-

- P1 10
- P2 1
- P3 0
- P4 2

Avg WT=13/4=3.25

Question 2(d):

Gantt chart:

P1	P2	P1	P3	P2	P1	P3	P2	
0	2	4	6	8	9	11	13	16

WT-

$$P1 \rightarrow (0-0)+(4-2)+(10-6)=6$$

$$P2 \rightarrow (2-1)+(8-4)+(13-10)=8$$

$$P3 \rightarrow (6-3)+(11-8)=6$$

$$\text{AVG WT:- } 20/3=6.67$$

TAT:-

$$P1:- (11-0)=11$$

$$P2:- (16-1)=15$$

$$P3:- (13-3)=10$$

$$\text{AVG TAT:- } 36/3=12$$

SECTION – C

Question 3(a): Batch System

- ❖ Batch means set of jobs with similar need.
- ❖ In this system, Operator makes Batch of jobs and loads it into punch card.
- ❖ Similar types of jobs batch together and execute at a time.

Advantages

- ✓ Save time from activities like loading compiler.
- ✓ No manual intervention by user is needed.

Disadvantages

- ✓ Limited Memory
- ✓ Interaction of I/O directly with CPU.
- ✓ CPU is often idle.

Question 3(b): Multiprogramming Operating System

- ❖ Main objective of multiprogramming is to maximize CPU utilization.
- ❖ More than one process can reside in main memory which are ready to execute i.e. more than one process is in ready state.
- ❖ In this system if any running process performs I/O or other event which does not require CPU, then instead of sitting idle, CPU performs context switching and picks another process for execution.

Advantages

- ✓ High CPU utilization.
- ✓ Less waiting time, response time etc.
- ✓ Useful in current scenario when load is high.

Disadvantages

- ✓ Process scheduling is difficult.
- ✓ Main memory management is required.
- ✓ Problems like memory fragmentation may occur.

Multiprocessing System

- ❖ A system is called multiprocessing system if two or more CPU within a single computer communicate with each other and share system bus memory and I/O devices.
- ❖ It provides true parallel execution of processes.

Advantages:

- ✓ **Increased Throughput:** By increasing number of processors, we expect to get more work done in less time.
- ✓ **Economy of Scale:** multiprocessor system can cost less than equivalent multiple single processor systems.
- ✓ **Increased Reliability:** Since work is distributed among several processors; failure of one processor will not halt the system only slow it down.

Disadvantages

- ✓ It is complex system.
- ✓ Process scheduling is difficult in this system.
- ✓ It required large size of Main memory.
- ✓ Overhead reduces throughput.

Multithreading

Thread: Light weight Process is called Thread.

- Light weight Process
- Thread Switching does not need interaction with OS
- Thread can share Code segment, Data Section, File Descriptor, address space, heap etc

Threads has its own

Register
Program Counter
Stack

Threads of same process share

Data Section
Code Segment
Heap
Address Space
File Descriptor
Message Queue

Note: Local variables of the Process are stored in Stack, Global variables are stored in Data Section and dynamically created variables are stored in Heap.

Question 4(a): Necessary conditions for deadlock

1. **Mutual Exclusion:** At least one resource must be held in a non sharable mode i.e. only one process at a time can use the resource. If another process requests the resource, the requesting process must be delayed until the resource has been released.
2. **Hold and wait:** A process must be holding one resource and waiting to acquire additional resources that are currently being held by other processes.
3. **No Preemption:** Resources cannot be preempted i.e. a resource can be released only voluntarily by the process holding it, after that process has completed its task.
4. **Circular wait:** A set $\{P_0, P_1, \dots, P_n\}$ of waiting processes must exist such that P_0 is waiting for a resource held by P_1 , P_1 is waiting for a resource held by P_2 , ... P_{n-1} is waiting for a resource held by P_n and P_n is waiting for a resource held by P_0 .

Question 4(b):

Need $[i, j] = \text{Max } [i, j] - \text{Allocation } [i, j]$

So, the content of Need Matrix is:

Process	Need		
	A	B	C
P ₀	7	4	3
P ₁	1	2	2
P ₂	6	0	0
P ₃	0	1	1
P ₄	4	3	1

Applying the Safety algorithm on the given system,

m=3, n=5 Step 1 of Safety Algo
 Work = Available
 Work =

3	3	2
---	---	---

0 1 2 3 4
 Finish =

false	false	false	false	false
-------	-------	-------	-------	-------

For i=3 Step 2
 Need₃ = 0, 1, 1 0, 1, 1 5, 3, 2
 Finish [3] = false and **Need₃ < Work**
 So P₃ must be kept in safe sequence

Step 3
 Work = Work + Allocation₀
 Work =

7	5	5
---	---	---

0 1 2 3 4
 Finish =

true	true	false	true	true
------	------	-------	------	------

For i=0 Step 2
 Need₀ = 7, 4, 3 7, 4, 3 3, 3, 2
 Finish [0] is false and **Need₀ > Work**
 So P₀ must wait But Need ≤ Work

Step 3
 Work = Work + Allocation₃
 Work =

7	4	3
---	---	---

0 1 2 3 4
 Finish =

false	true	false	true	false
-------	------	-------	------	-------

For i=2 Step 2
 Need₂ = 6, 0, 0 6, 0, 0 7, 5, 5
 Finish [2] is false and **Need₂ < Work**
 So P₂ must be kept in safe sequence

For i=1 Step 2
 Need₁ = 1, 2, 2 1, 2, 2 3, 3, 2
 Finish [1] is false and **Need₁ < Work**
 So P₁ must be kept in safe sequence

For i=4 Step 2
 Need₄ = 4, 3, 1 4, 3, 1 7, 4, 3
 Finish [4] = false and **Need₄ < Work**
 So P₄ must be kept in safe sequence

Step 3
 Work = Work + Allocation₂
 Work =

10	5	7
----	---	---

0 1 2 3 4
 Finish =

true	true	true	true	true
------	------	------	------	------

Step 3
 Work = Work + Allocation₁
 Work =

5	3	2
---	---	---

0 1 2 3 4
 Finish =

false	true	false	false	false
-------	------	-------	-------	-------

Step 3
 Work = Work + Allocation₄
 Work =

7	4	5
---	---	---

0 1 2 3 4
 Finish =

false	true	false	true	true
-------	------	-------	------	------

Step 4
 Finish [i] = true for 0 ≤ i ≤ n
 Hence the system is in Safe state

For i=2 Step 2
 Need₂ = 6, 0, 0 6, 0, 0 5, 3, 2
 Finish [2] is false and **Need₂ > Work**
 So P₂ must wait

For i=0 Step 2
 Need₀ = 7, 4, 3 7, 4, 3 7, 4, 5
 Finish [0] is false and **Need < Work**
 So P₀ must be kept in safe sequence

The safe sequence is P₁, P₃, P₄, P₀, P₂

Question 5(a):

Operating System

- ❖ Operating system is a system program that act as an interface between hardware and user.
- ❖ It manages system resources.
- ❖ It provides a platform on which other application programs are installed.
- ❖ Core Part of Operating System is called kernel.

Goals of operating System

Primary Goal:To make system user friendly (convenient).

Secondary Goal:To make system efficient.

Role of Operating System: We can explore operating Systems from two viewpoints that of the user and that of the system.

User's View: In user's view operating system is designed to maximize the work (CPU utilization) that the user is performing. In this context OS is designed mostly for **ease of use** with some attention paid to performance and none paid to resource utilization.

System's View: In system point of view operating system is a program that works as a **resource allocator**. Every computer system has many resources like CPU Time, Memory Space, File Storage Space, I/O Device and so on. The operating system act as manager of these resources.

Functions of Operating System

Operating System is responsible for following activities.

Process Management: Operating System is responsible for following activities in connection with Process Management.

- ✓ Scheduling Processes and threads on the CPUs.
- ✓ Creating and deleting both user and system processes.
- ✓ Suspending and resuming processes.
- ✓ Providing mechanisms for process synchronization.
- ✓ Providing mechanism for process communication.

Memory Management: Operating System is responsible for following activities in connection with Memory Management.

- ✓ Keeping track of which parts of memory are currently being used and by whom.
- ✓ Deciding which processes and data to move into and out of memory.
- ✓ Allocating and de-allocating memory space as needed.

Disk Management: Operating System is responsible for following activities in connection with Disk Management.

- ✓ Free space management in disk.
- ✓ Storage Allocation in disk.
- ✓ Disk Scheduling.

File Management: Operating System is responsible for following activities in connection with File Management.

- ✓ Creating and deleting files.
- ✓ Creating and deleting directories to organize files.
- ✓ Supporting primitives for manipulating files and directories.
- ✓ Mapping files onto secondary storage.
- ✓ Backing up files on stable storage media.

I/O Device Management: Operating System is responsible for I/O Device management needed by various processes.

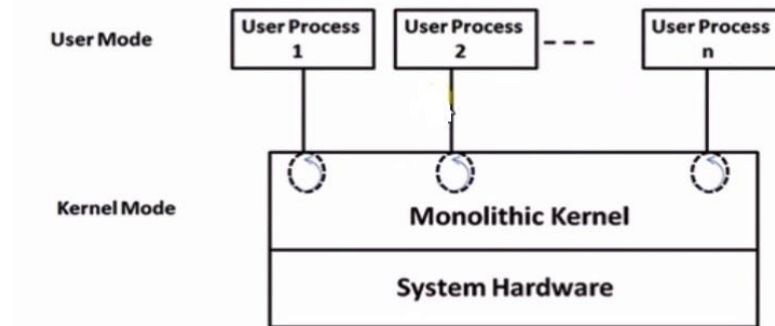
Network Management Operating System is responsible for network management.

Security and Protection: Operating System also provides security and protection to computer system.

Question 5(b):

Monolithic Kernel

- ❖ Monolithic means all in one piece.
 - ❖ In monolithic kernel user services and kernel services are implemented under same address space, It increases size of the kernel thus increase size of operating system.
- Example: Linux



Advantage

- ✓ Execution is faster.

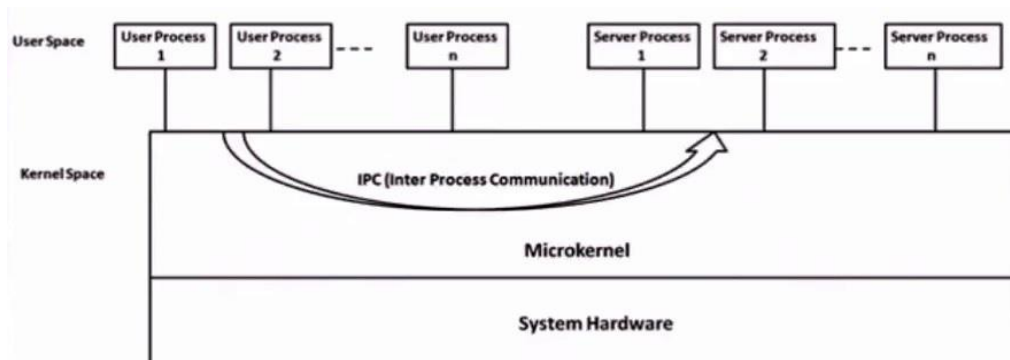
Disadvantages

- ✓ Size is larger.
- ✓ Hard to extend.
- ✓ Hard to port.
- ✓ More prone to errors and bugs.

Microkernel

- ❖ In microkernel, user services and kernel services are implemented in different address space, therefore it also reduces the size of kernel as well as the size of Operating system.
- ❖ In this architecture, only the most important services are present inside kernel and rest of the operating system services are present inside system application program.
- ❖ Communication between client process and services running in user address space is established through message passing thus reduce the speed of execution.

Example: Mach OS



Advantages

- ✓ Size is smaller.
- ✓ Easy to extend.
- ✓ Easy to port.
- ✓ Less prone to errors and bugs.

Disadvantage

- ✓ Execution is slower.

Difference between microkernel and monolithic kernel

Parameter	Microkernel	Monolithic Kernel
Size	Small	Large
Speed of execution	Faster	Slower
Extendibility	Easy	Hard
Portability	Easy	Hard
Error prone capability	Less	More